



# **L(2,1) Graph Labelling**

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# Overview

1

What's L(2,1) labeling?

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2

Mathematical  
Bounds

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3

Our  
Implementation

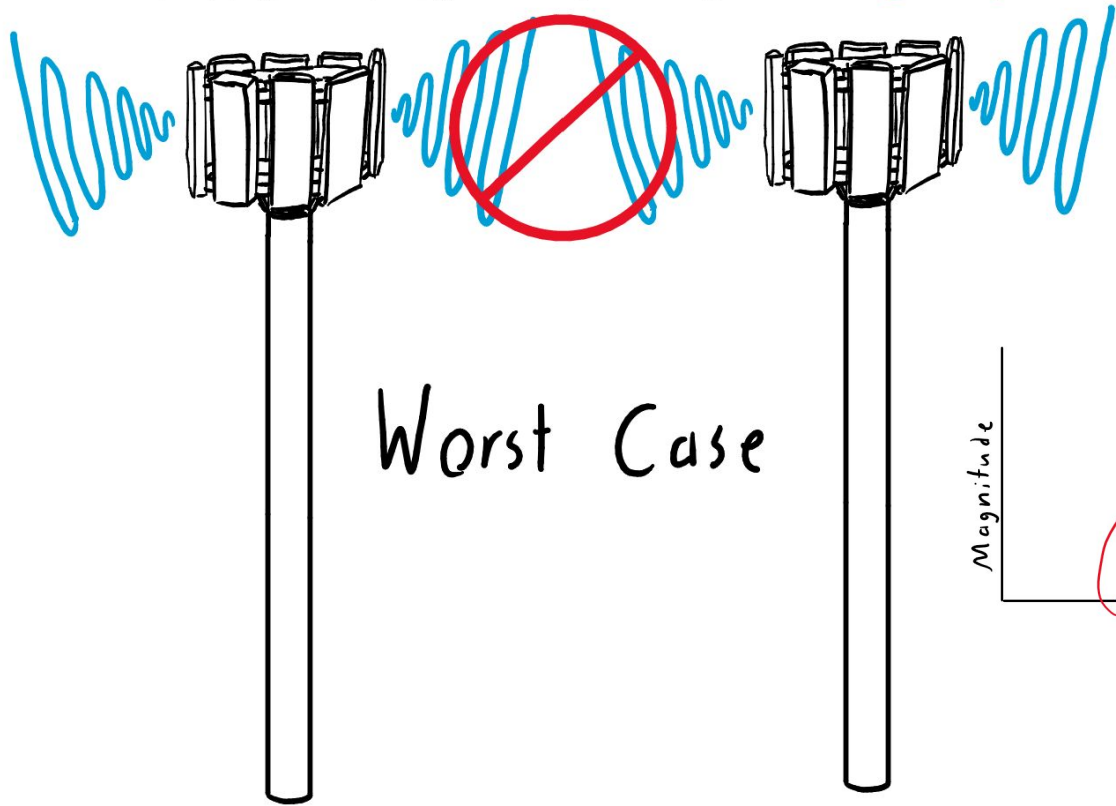


# What's $L(2,1)$ Labeling?

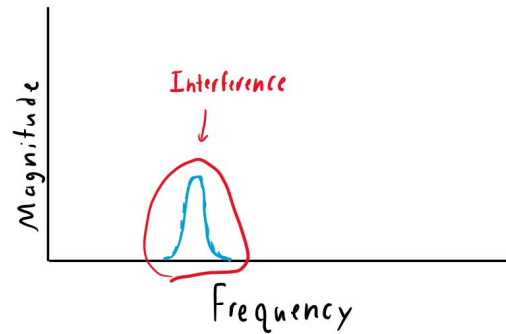
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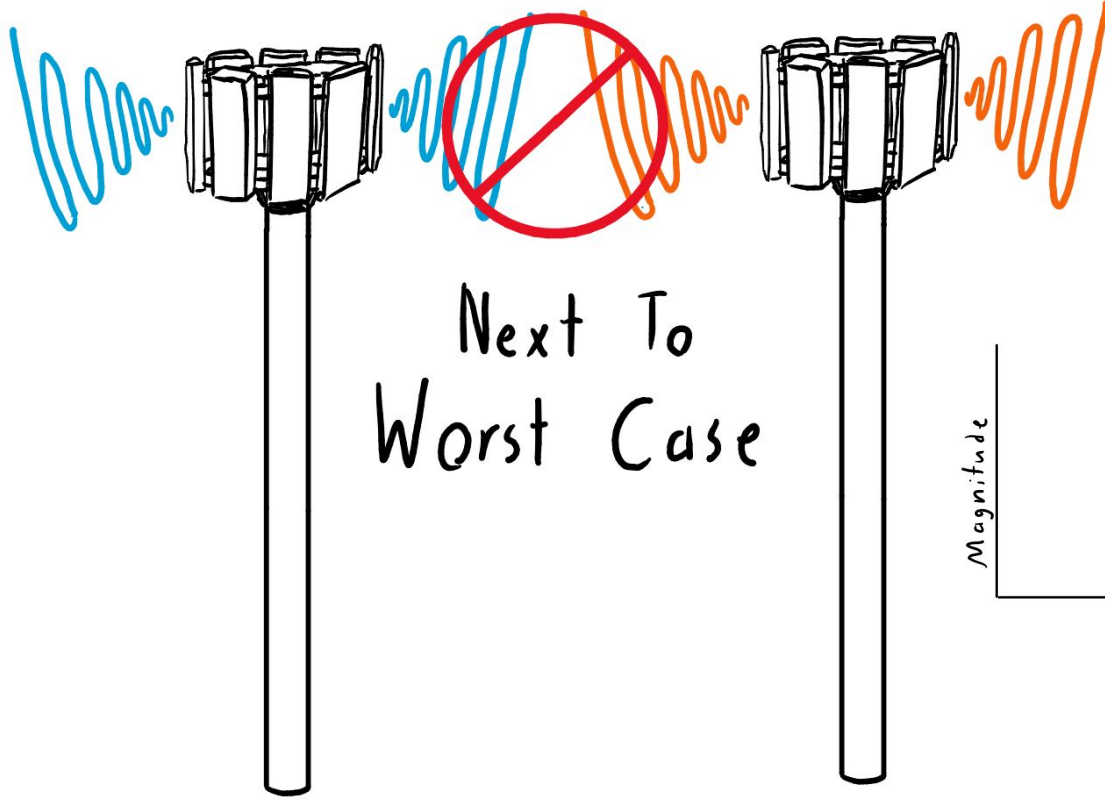
# Two Towers On Band Blue



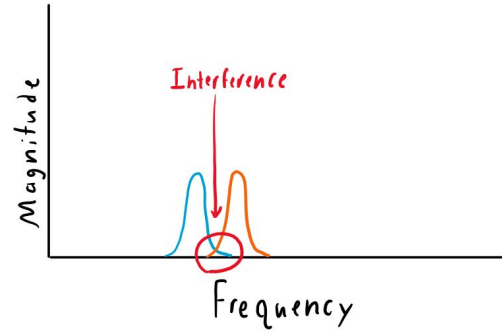
Worst Case




# Two Towers On Adjacent Bands



Next To Worst Case






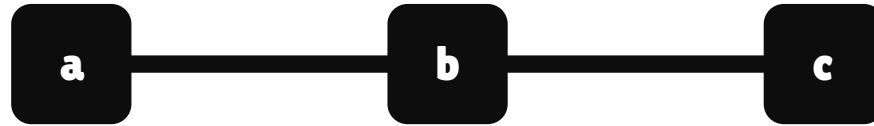
# L(2,1)–Labeling: A special graph coloring

For a graph, whenever  $x$  and  $y$  are two adjacent vertices then their label must have a distance greater than or equal to two.

Whenever  $x$  and  $y$  are two vertices with distance **two** between them, then their label must have a distance greater than or equal to **one**.

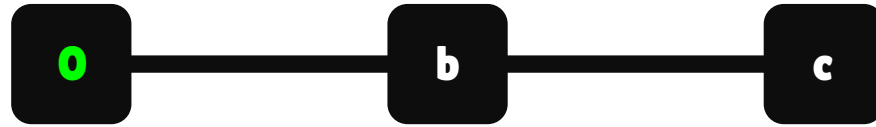


**Let's do an example!**



**Say we have a three vertex, two edge graph, and we want to assign them a minimum label from the set  $\{0,1,2,3,4,\dots\}$**

**Let's do an example!**





**Let's do an example!**



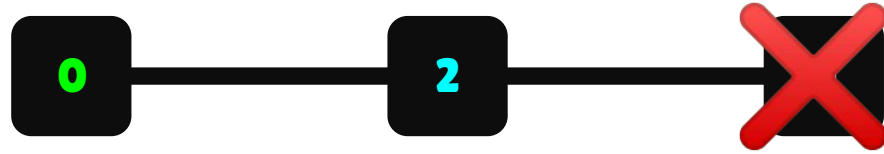
**Let's do an example!**



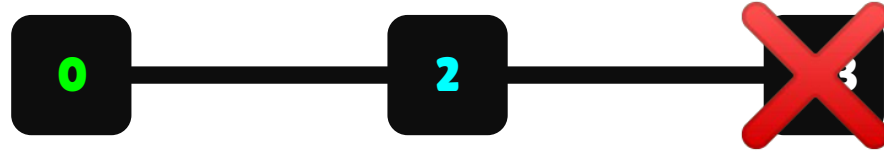
**Let's do an example!**



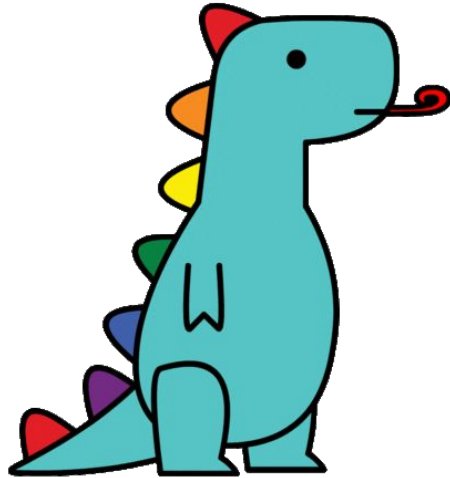
**Let's do an example!**



**Let's do an example!**



**Done! But, we do have two holes...**



# **Mathematical Bounds: Minimizing the largest label**



# Finding $\lambda(G)$ : The minimum span of labels required $\{0, 1, \dots, \lambda\}$

Griggs and Yeh's Conjectured Bound for  $\lambda(G)$

$$\lambda(G) \leq \Delta^2$$

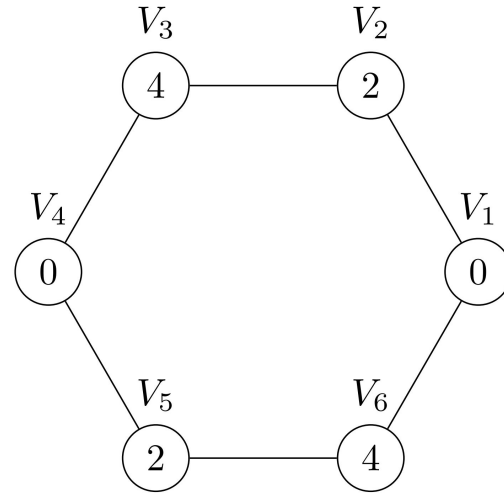
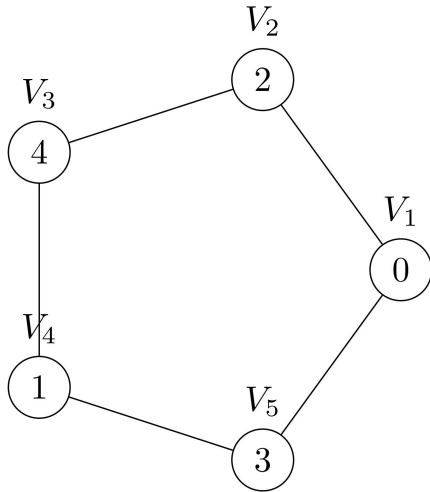
$\Delta$  = Highest degree vertex in a graph  $G$



$$\Delta^2 = 4$$



# Upper Bound for Cycles

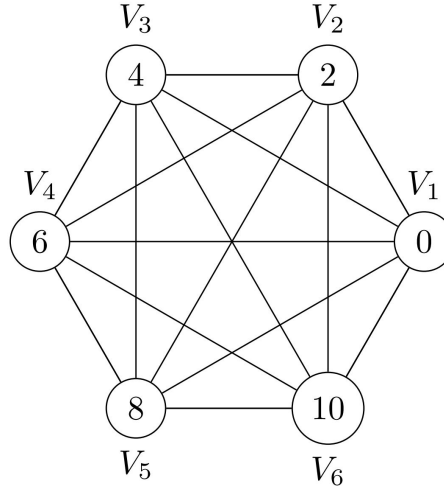


$\lambda = 4$  for  $C(n)$  where  $n \geq 3$

Some cycles have holes; others do not

$\Delta = 2$ , so cycles meet the conjectured  $\Delta^2$  bound

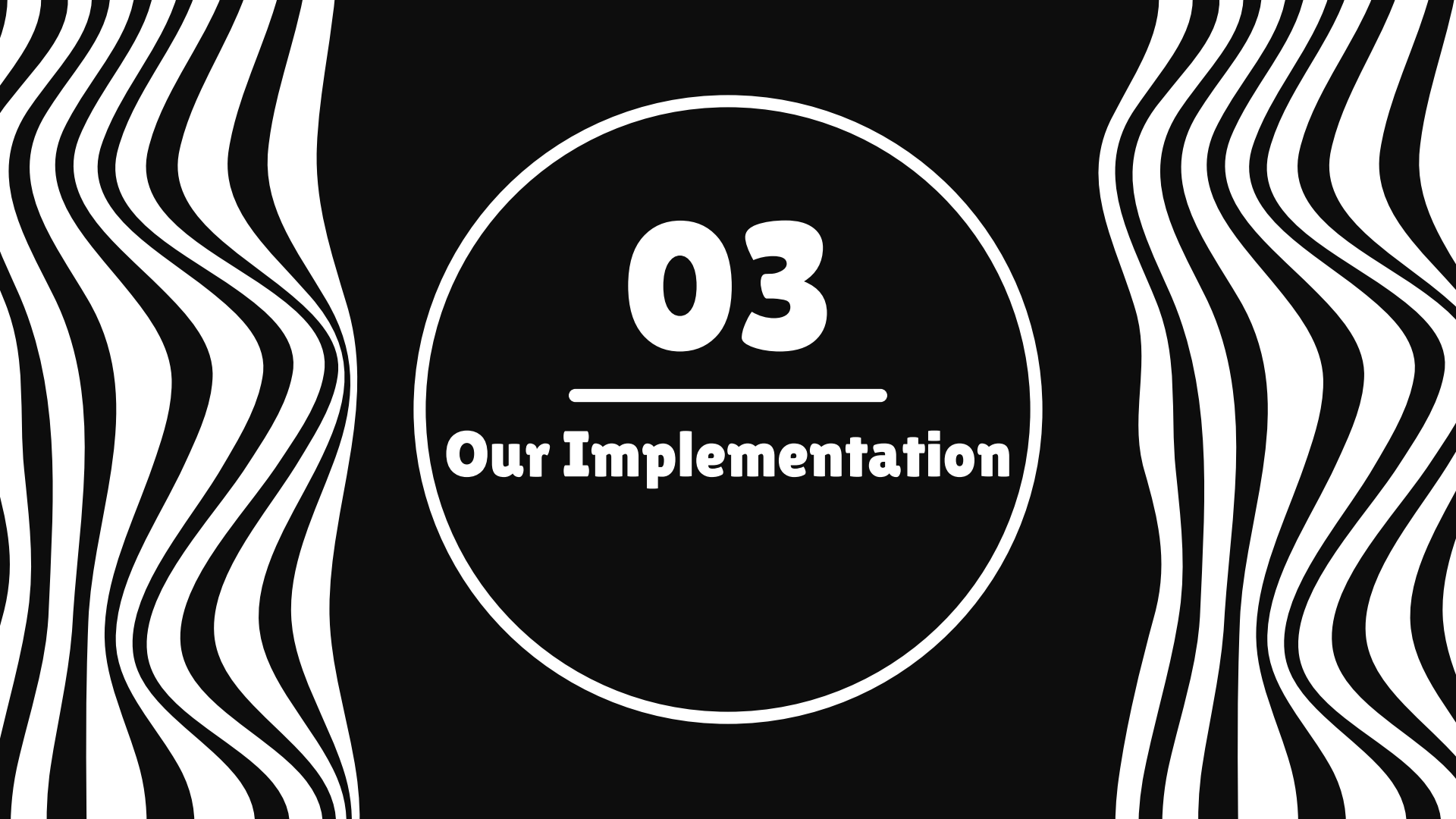
# Upper Bound for Complete Graphs



$\lambda = 2(n-1)$  for  $K(n)$  where  $n \geq 3$

Only even labeling numbers are used

$\Delta = n-1$ , so complete graphs meet the conjectured  $\Delta^2$  bound



**03**

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**Our Implementation**

# Algorithms for Minimizing $\lambda(G)$

$\Delta$  = Highest degree vertex in a graph  $G$

**1**

Greedy Algorithm

$$\Delta^2 + 2\Delta$$

**2**

Modified Chang-Kuo Algorithm

$$\Delta^2 + \Delta - 2$$

**3**

Griggs and Yeh's Conjectured Bound  
(No known algorithm)

$$\Delta^2$$



# Algorithmic Differences

## Greedy Algorithm

- Iterate through vertices
- Assign lowest possible number
- Not seeing bigger picture, so prone to holes

## Modified Chang-Kuo Algorithm

- Iterate through labeling numbers
- Assign current number to as many vertices as possible
- Looks at entire graph each time, so reduces likelihood of holes

# Our approach

**1**

Takes a list of edge connections defined by the user

**2**

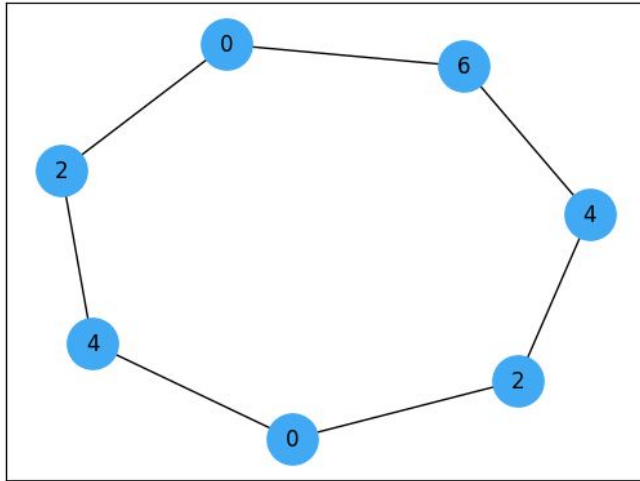
Converts this into an adjacency matrix

**3**

Solves and plots the labeled graphs with either the greedy or modified Chang-Kuo algorithm

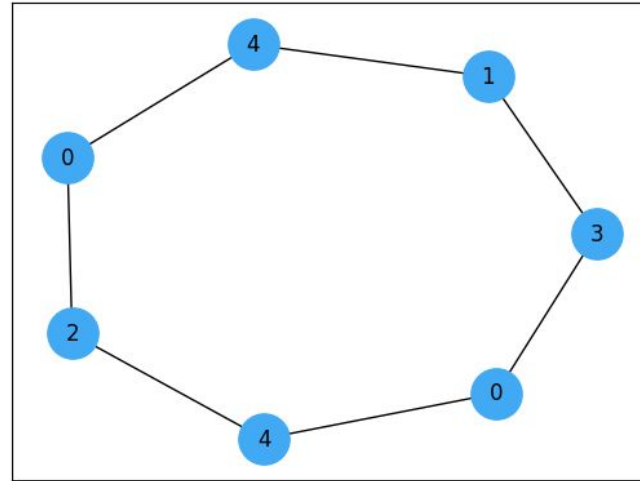
# Our Results (Cyclical Graph)

Greedy



$\lambda = 6$ , Holes = 3

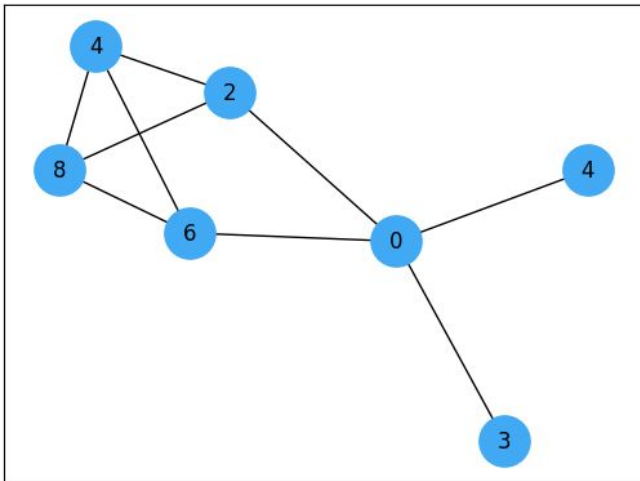
Chang-Kuo



$\lambda = 4$ , Holes = 0

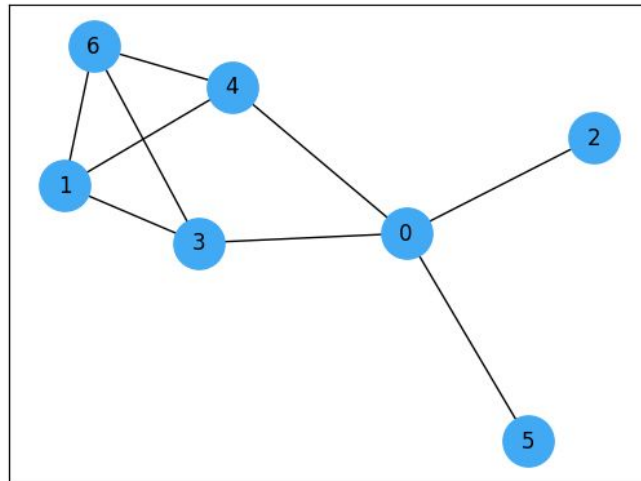
# Our Results (Random Graph)

Greedy



$\lambda = 8$ , Holes = 3

Chang-Kuo



$\lambda = 6$ , Holes = 0



# Check out our Code!

<https://github.com/olincollage/L21-Graph-Coloring/tree/main>





**Thank You!**  
**Special Thanks to Professor**  
**Nathaniel Karst of Babson (Olin '07)**  
**And Sarah Spence Adams!**

